

S6.D. Mixing Flow

A system is called **mixing** if

$$\begin{aligned} & \lim_{t \rightarrow \pm\infty} \frac{1}{\Sigma(E)} \int_{S_E} dS_E f(\mathbf{X}^N) g(\mathbf{X}^N(t)) \\ &= \frac{1}{[\Sigma(E)]^2} \int_{S_E} dS_E f(\mathbf{X}^N) \int_{S_E} dS_E g(\mathbf{X}^N(t)) \quad \forall f, g \in L^2 \end{aligned} \quad (6.113)$$

where $\Sigma(E)$ is the structure function defined in §6.C.

Setting $g(\mathbf{X}^N(t)) = \rho(\mathbf{X}^N, t)$ with

$$\int_{S_E} dS_E \rho(\mathbf{X}^N, t) = \Sigma(E) \quad \forall t$$

then, for a mixing system,

$$\begin{aligned} \langle f(t) \rangle &= \lim_{t \rightarrow \pm\infty} \frac{1}{\Sigma(E)} \int_{S_E} dS_E f(\mathbf{X}^N) \rho(\mathbf{X}^N, t) \\ &= \frac{1}{[\Sigma(E)]} \int_{S_E} dS_E f(\mathbf{X}^N) \\ &= \int_{S_E} dS_E f(\mathbf{X}^N) \rho_s \end{aligned} \quad (6.114)$$

where

$$\rho_s = \frac{1}{\Sigma(E)} = \rho(\mathbf{X}^N, t \rightarrow \pm\infty)$$

Thus, a mixing system behaves like a microcanonical ensemble as $t \rightarrow \pm\infty$ [see (6.46)].

Note however that owing to the balance equation [see (6.23-4)]

$$\frac{d\rho(\mathbf{X}^N, t)}{dt} = \left(\frac{\partial}{\partial t} + \dot{\mathbf{X}}^N \cdot \nabla_{\mathbf{X}^N} \right) \rho = 0$$

the phase space volume occupied by a given set of phase points remains unchanged in the course of dynamic evolution. Therefore, starting with a non-equilibrium ρ that is non-zero only over a portion of the energy surface S_E , it can never evolve into ρ_s that is non-zero everywhere on S_E . The best one can hope for is a coarse-grained uniformity such that the average value of ρ over any small region of some fixed size is the same everywhere on S_E . This is akin to the mixing of oil and water as shown in figure 6.2 in Reichl's text.

Baker's Map

A binary number

$$r = (0.c_{-1}c_{-2}c_{-3}\dots)_2 \quad \text{with } c_k = 0 \text{ or } 1 \text{ for } k = -1, -2, -3, \dots$$

corresponds to a (decimal) number $r \in [0, 1]$ of value

$$r = \sum_{k=-1}^{-\infty} c_k 2^k = \sum_{k=1}^{\infty} c_{-k} 2^{-k} = \frac{c_{-1}}{2} + \frac{c_{-2}}{4} + \frac{c_{-3}}{8} + \dots \quad (6.115a)$$

Note that

$$(0.000\dots)_2 = 0$$

$$(0.111 \dots)_2 = \sum_{k=1}^{\infty} 2^{-k} = \frac{\frac{1}{2} - \left(\frac{1}{2}\right)^{\infty}}{1 - \frac{1}{2}} = 1$$

$$(0.1000 \dots)_2 = \frac{1}{2}$$

A point within a unit square has coordinates

$$(p, q) \quad \text{with} \quad 0 \leq p \leq 1 \text{ and } 0 \leq q \leq 1$$

Let the binary form of (p, q) be

$$(p, q) = (0.p_{-1} p_{-2} p_{-3} \dots, 0.q_{-1} q_{-2} q_{-3} \dots)_2 \quad (6.115b)$$

where

$$p_k, q_k = 0 \text{ or } 1 \quad \text{for} \quad k = -1, -2, -3, \dots$$

We can represent it as a real number $s \in (-\infty, 0]$ with

$$\begin{aligned} s &= (\dots s_2 s_1 s_0 \cdot s_{-1} s_{-2} s_{-3} \dots)_2 \\ &= (\dots p_{-3} p_{-2} p_{-1} \cdot q_{-1} q_{-2} q_{-3} \dots)_2 \\ &= \sum_{k=-\infty}^{\infty} s_k 2^k \end{aligned} \quad (6.115c)$$

Thus,

$$\begin{aligned} s_k &= \begin{cases} p_{-(k+1)} & \text{for } k = 0, 1, 2, \dots \\ q_k & \text{for } k = -1, -2, \dots \end{cases} \\ p &= \sum_{k=-1}^{\infty} p_k 2^k = \sum_{k=0}^{\infty} s_k 2^{-(k+1)} \\ q &= \sum_{k=-1}^{\infty} q_k 2^k = \sum_{k=-1}^{\infty} s_k 2^k \end{aligned} \quad (6.115)$$

The right bit-shift (or Bernoulli shift) operation shifts all digits of a binary number to the right by one position, e.g.,

$$\begin{aligned} R s_k &= s_{k+1} \\ R s &= (\dots s_3 s_2 s_1 \cdot s_0 s_{-1} s_{-2} \dots)_2 \\ &= (\dots p_{-4} p_{-3} p_{-2} \cdot p_{-1} q_{-1} q_{-2} \dots)_2 \\ &= \sum_{k=-\infty}^{\infty} s_{k+1} 2^k = \sum_{k=-\infty}^{\infty} s_k 2^{k-1} = \frac{1}{2} s \\ &= s' = (\dots s'_2 s'_1 s'_0 \cdot s'_{-1} s'_{-2} s'_{-3} \dots)_2 \end{aligned}$$

The right bit-shift $s' = R s$ corresponds to a transformation U on (p, q) given by

$$\begin{aligned} (p', q') &= U(p, q) = R s \\ &= (\dots p'_{-3} p'_{-2} p'_{-1} \cdot q'_{-1} q'_{-2} q'_{-3} \dots)_2 \\ \rightarrow \quad p' &= (0.p_{-2} p_{-3} p_{-4} \dots)_2 = \sum_{k=-2}^{\infty} p_k 2^{k+1} = -p_{-1} + \sum_{k=-1}^{\infty} p_k 2^{k+1} = -p_{-1} + 2p \\ q' &= (0.p_{-1} q_{-1} q_{-2} \dots)_2 = \frac{1}{2} p_{-1} + \sum_{k=-1}^{\infty} q_k 2^{k-1} = \frac{1}{2} p_{-1} + \frac{1}{2} q \end{aligned}$$

Hence,

$$\begin{aligned}
U(p, q) &= \left(-p_{-1} + 2p, \frac{1}{2}p_{-1} + \frac{1}{2}q \right) \\
&= \begin{cases} \left(2p, \frac{1}{2}q \right) & \text{for } p_{-1} = 0 \text{ or } 0 \leq p \leq \frac{1}{2} \\ \left(-1 + 2p, \frac{1}{2} + \frac{1}{2}q \right) & \text{for } p_{-1} = 1 \text{ or } \frac{1}{2} \leq p \leq 1 \end{cases} \quad (6.116)
\end{aligned}$$

The inverse of the right bit-shift is the left bit-shift operation which shifts all digits of a binary number to the left by one position, i.e.,

$$\begin{aligned}
L s_k &= s_{k-1} \\
L s &= (\dots s_2 s_1 s_0 s_{-1} \cdot s_{-2} s_{-3} \dots)_2 \\
&= (\dots p_{-3} p_{-2} p_{-1} q_{-1} \cdot q_{-2} q_{-3} \dots)_2 \\
&= \sum_{k=-\infty}^{\infty} s_{k-1} 2^k = \sum_{k=-\infty}^{\infty} s_k 2^{k+1} = 2s \\
&= s' = (\dots s'_2 s'_1 s'_0 \cdot s'_{-1} s'_{-2} s'_{-3} \dots)_2
\end{aligned}$$

The left bit-shift $s' = Ls$ corresponds to a transformation U^{-1} on (p, q) given by

$$\begin{aligned}
(p', q') &= U^{-1}(p, q) = Ls \\
&= (\dots p'_{-3} p'_{-2} p'_{-1} \cdot q'_{-1} q'_{-2} q'_{-3} \dots)_2 \\
\rightarrow p' &= (0. q_{-1} p_{-1} p_{-2} \dots)_2 = \frac{q_{-1}}{2} + \sum_{k=-1}^{-\infty} p_k 2^{k-1} = \frac{q_{-1}}{2} + \frac{1}{2}p \\
q' &= (0. q_{-2} q_{-3} q_{-4} \dots)_2 = \sum_{k=-2}^{-\infty} q_k 2^{k+1} = -q_{-1} + \sum_{k=-1}^{-\infty} q_k 2^{k+1} = -q_{-1} + 2q
\end{aligned}$$

Hence,

$$\begin{aligned}
U^{-1}(p, q) &= \left(\frac{q_{-1}}{2} + \frac{1}{2}p, -q_{-1} + 2q \right) \\
&= \begin{cases} \left(\frac{1}{2}p, 2q \right) & \text{for } q_{-1} = 0 \text{ or } 0 \leq q \leq \frac{1}{2} \\ \left(\frac{1}{2} + \frac{1}{2}p, -1 + 2q \right) & \text{for } q_{-1} = 1 \text{ or } \frac{1}{2} \leq q \leq 1 \end{cases} \quad (6.117)
\end{aligned}$$

The Jacobian of U is

$$J = \begin{vmatrix} \frac{\partial p'}{\partial p} & \frac{\partial p'}{\partial q} \\ \frac{\partial q'}{\partial p} & \frac{\partial q'}{\partial q} \end{vmatrix} = \begin{vmatrix} 2 & 0 \\ 0 & \frac{1}{2} \end{vmatrix} = 1$$

so that the map is area preserving. So is the inverse map U^{-1} .

For a given probability density $\rho(p, q)$ on the unit square, U induces a (time) evolution on ρ given by

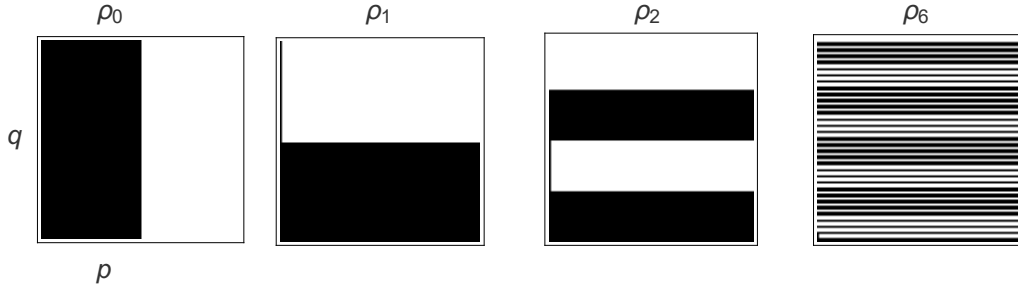
$$\begin{aligned}
\rho(p, q, t_n) &\equiv \rho_n(p, q) = U^n \rho_0(p, q) = U^n \rho(p, q, 0) \\
&= \rho_0(U^{-n} p, U^{-n} q) \quad (6.118)
\end{aligned}$$

In particular,

$$\rho_{n+1}(p, q) = U \rho_n(p, q) = \rho_n(U^{-1} p, U^{-1} q)$$

$$= \begin{cases} \rho_n \left(\frac{1}{2} p, 2q \right) & \text{for } 0 \leq q \leq \frac{1}{2} \\ \rho_n \left(\frac{1}{2} + \frac{1}{2} p, -1 + 2q \right) & \text{for } \frac{1}{2} \leq q \leq 1 \end{cases} \quad (6.119)$$

The following are results obtained from the *Mathematica* code in §Code. [Fig.6.3. in Reichl's text was printed upside-down.]



Note that the initial ρ_0 , which occupies only the left side of the square, is flung down after one mapping to the bottom half of the square. Subsequent mappings are like stretching the non-zero area and then folding it back on top of itself, which is reminiscent of a baker kneading dough. Hence the name of the map.

Since $s_0 = p_{-1}$, therefore,

$$\begin{aligned} p = (0.0 p_{-2} p_{-3} \dots)_2 \leq (0.1)_2 = \frac{1}{2} & \quad \text{if } s_0 = 0 \\ p = (0.1 p_{-2} p_{-3} \dots)_2 \geq (0.1)_2 = \frac{1}{2} & \quad \text{if } s_0 = 1 \end{aligned}$$

i.e., p lies to the left (right) of $\frac{1}{2}$ for $s_0 = 0 (1)$.

Consider two nearby points s and s' with

$$s_k = s'_k \quad \forall k = -m, -m + 1, \dots, \infty$$

They will remain close-by for the first m left bit-shifts until s_0 and s'_0 is replaced by s_{-m-1} and s'_{-m-1} , respectively. These points are then flung to the opposite sides of $p = \frac{1}{2}$. The mapping is therefore

chaotic.

Consider now the reduced probability defined by

$$\phi(p) = \int_0^1 dq \rho(p, q)$$

and the corresponding evolutions

$$\phi_n(p) = \int_0^1 dq \rho_n(p, q)$$

Using (6.119) we have

$$\begin{aligned} \phi_{n+1}(p) &= \int_0^{1/2} dq \rho_n \left(\frac{1}{2} p, 2q \right) + \int_{1/2}^1 dq \rho_n \left(\frac{1}{2} p + \frac{1}{2}, 2q - 1 \right) \\ &= \frac{1}{2} \int_0^1 dq' \rho_n \left(\frac{1}{2} p, q' \right) + \int_0^1 dq'' \rho_n \left(\frac{1}{2} p + \frac{1}{2}, q'' \right) \end{aligned}$$

where

$$q' = 2q \rightarrow dq = \frac{1}{2} dq' \quad \text{and} \quad q' \in [0, 1] \quad \text{for} \quad q \in \left[0, \frac{1}{2}\right]$$

$$q'' = 2q - 1 \rightarrow dq = \frac{1}{2} dq'' \quad \text{and} \quad q'' \in [0, 1] \quad \text{for} \quad q \in \left[\frac{1}{2}, 1\right]$$

Therefore,

$$\phi_{n+1}(\rho) = \frac{1}{2} \phi_n\left(\frac{1}{2}\rho\right) + \frac{1}{2} \phi_n\left(\frac{1}{2}\rho + \frac{1}{2}\right) \quad (6.120)$$

which is a Markovian chain [c.f.(5.20)].

Starting at $n = 0$, we have

$$\begin{aligned} \phi_1(\rho) &= \frac{1}{2} \left[\phi_0\left(\frac{1}{2}\rho\right) + \phi_0\left(\frac{1}{2}\rho + \frac{1}{2}\right) \right] \\ \phi_2(\rho) &= \frac{1}{2} \left[\phi_1\left(\frac{1}{2}\rho\right) + \phi_1\left(\frac{1}{2}\rho + \frac{1}{2}\right) \right] \\ &= \frac{1}{2^2} \left[\phi_0\left(\frac{1}{2^2}\rho\right) + \phi_0\left(\frac{1}{2^2}\rho + \frac{1}{2}\right) + \phi_0\left(\frac{1}{2^2}\rho + \frac{1}{2^2}\right) + \phi_0\left(\frac{1}{2^2}\rho + \frac{3}{2^2}\right) \right] \\ &= \frac{1}{2^2} \sum_{k=0}^3 \phi_0\left(\frac{1}{2^2}\rho + \frac{k}{2^2}\right) \\ \phi_3(\rho) &= \frac{1}{2} \left[\phi_2\left(\frac{1}{2}\rho\right) + \phi_2\left(\frac{1}{2}\rho + \frac{1}{2}\right) \right] \\ &= \frac{1}{2^3} \sum_{k=0}^3 \left[\phi_0\left(\frac{1}{2^3}\rho + \frac{k}{2^2}\right) + \phi_0\left(\frac{1}{2^3}\rho + \frac{1}{2^3} + \frac{k}{2^2}\right) \right] \\ &= \frac{1}{2^3} \sum_{k=0}^3 \left[\phi_0\left(\frac{1}{2^3}\rho + \frac{2k}{2^3}\right) + \phi_0\left(\frac{1}{2^3}\rho + \frac{2k+1}{2^3}\right) \right] \\ &= \frac{1}{2^3} \sum_{k=0}^7 \phi_0\left(\frac{1}{2^3}\rho + \frac{k}{2^3}\right) \\ \rightarrow \phi_n(\rho) &= \frac{1}{2^n} \sum_{k=0}^{2^n-1} \phi_0\left(\frac{1}{2^n}\rho + \frac{k}{2^n}\right) \quad (6.121) \end{aligned}$$

Proof by induction:

Using (6.120), we have

$$\begin{aligned} \phi_{n+1}(\rho) &= \frac{1}{2} \phi_n\left(\frac{1}{2}\rho\right) + \frac{1}{2} \phi_n\left(\frac{1}{2}\rho + \frac{1}{2}\right) \\ &= \frac{1}{2^{n+1}} \sum_{k=0}^{2^n-1} \left[\phi_0\left(\frac{1}{2^{n+1}}\rho + \frac{k}{2^n}\right) + \phi_0\left(\frac{1}{2^{n+1}}\rho + \frac{1}{2^{n+1}} + \frac{k}{2^n}\right) \right] \\ &= \frac{1}{2^{n+1}} \sum_{k=0}^{2^n-1} \left[\phi_0\left(\frac{1}{2^{n+1}}\rho + \frac{2k}{2^{n+1}}\right) + \phi_0\left(\frac{1}{2^{n+1}}\rho + \frac{2k+1}{2^{n+1}}\right) \right] \\ &= \frac{1}{2^{n+1}} \sum_{k=0}^{2^{n+1}-1} \phi_0\left(\frac{1}{2^{n+1}}\rho + \frac{k}{2^{n+1}}\right) \quad [2(2^n - 1) + 1 = 2^{n+1} - 1] \end{aligned}$$

QED.

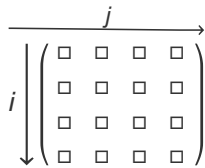
For any smooth function ϕ_0 , (6.121) gives

$$\begin{aligned}
 \lim_{n \rightarrow \infty} \phi_n(\rho) &= \lim_{n \rightarrow \infty} \frac{1}{2^n} \sum_{k=0}^{2^n-1} \phi_0\left(\frac{1}{2^n} \rho + \frac{k}{2^n}\right) \\
 &= \lim_{n \rightarrow \infty} \sum_{k=0}^{2^n-1} \frac{\Delta k}{2^n} \phi_0\left(\frac{1}{2^n} \rho + \frac{k}{2^n}\right) && \Delta k = 1 \\
 &= \lim_{n \rightarrow \infty} \int_0^{1-2^{-n}} dy \phi_0\left(\frac{1}{2^n} \rho + y\right) && y = \frac{k}{2^n} \\
 &= \int_0^1 dy \phi_0(y) \\
 &= 1 && (6.122)
 \end{aligned}$$

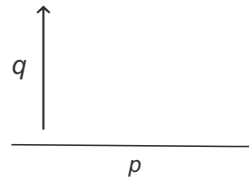
i.e., the reduced, or coarse-grained, probability density evolves to an equilibrium value, but not so for $n \rightarrow -\infty$. Thus, the Baker's transformation exhibits irreversibility in a coarse-grained sense.

Code

Labelling of array:



Labelling of graph:



(* This perform m Baker's transformations on the n×n matrix $\rho\theta$ *)

```

Baker[ $\rho\theta$ _, m_, n_] := Module[{ $\rho\theta\theta = \rho\theta$ ,  $\rho = \rho\theta$ },
  Do[
    Do[
       $p = \frac{n-j}{n-1}$ ;  $q = \frac{i-1}{n-1}$ ;
      If[ $q \leq \frac{1}{2}$ ,  $p\theta = \frac{1}{2} p$ ;  $q\theta = 2 q$ ,  $p\theta = \frac{1}{2} (p+1)$ ;  $q\theta = 2 q - 1$ ];
       $j\theta = \text{Round}[n - p\theta (n-1)]$ ;  $i\theta = \text{Round}[q\theta (n-1) + 1]$ ;
       $\rho\theta\theta[[i, j]] = \rho\theta\theta[[i\theta, j\theta]]$ ,
      {i, n}, {j, n}];
     $\rho\theta\theta = \rho$ ,
    m];
   $\rho$ ]

n = 101;
 $\rho\theta = \text{Table}[\theta, \{i, n\}, \{j, n\}]$ ;
Table[ $\rho\theta[[i // \text{Round}, j // \text{Round}]] = 1$ , {j, 1,  $\frac{n}{2} + 1$ }, {i, n}];

ArrayPlot[ $\rho\theta$ ]

 $\rho = \text{Baker}[\rho\theta, 1, n]$ ;
ArrayPlot[ $\rho$ ]

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 $\rho = \text{Baker}[\rho\theta, 2, n];$   
 $\text{ArrayPlot}[\rho]$ 
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 $\rho = \text{Baker}[\rho\theta, 6, n];$   
 $\text{ArrayPlot}[\rho]$ 
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Exercise 6.8.

Compute the trace of the Baker's map, \hat{U}^n .

Answer

Consider \hat{U} as an operator on the Hilbert space spanned by the basis vectors $|p, q\rangle$.

The map

$$\hat{U}^n(p, q) = (p^n, q^n)$$

can be written as

$$\hat{U}^n |p, q\rangle = |p^n, q^n\rangle$$

Hence,

$$\begin{aligned} \text{Tr } \hat{U}^n &= \int_0^1 dp \int_0^1 dq \langle p, q | \hat{U}^n | p, q \rangle \\ &= \int_0^1 dp \int_0^1 dq \langle p, q | p^n, q^n \rangle \\ &= \int_0^1 dp \int_0^1 dq \delta(p - p^n) \delta(q - q^n) \end{aligned} \quad (1)$$

Using

$$\delta(f(x)) = \sum_{k=1}^M \frac{\delta(x - x_k)}{|f'(x_k)|} \quad (2)$$

where

$$f(x_k) = 0 \quad \text{for } k = 1, \dots, M \quad f'(x) = \frac{df(x)}{dx}$$

we have

$$\begin{aligned} \delta(p - p^n) \delta(q - q^n) &= \sum_{k=1}^{M_n} \left| \frac{\partial p^n}{\partial p} - 1 \right|_{p=p_k, q=q_k}^{-1} \left| \frac{\partial q^n}{\partial q} - 1 \right|_{p=p_k, q=q_k}^{-1} \\ &\quad * \delta(q - q_k) \delta(p - p_k) \end{aligned} \quad (2a)$$

where (p_k, q_k) are points periodic under \hat{U} with period n , i.e.,

$$\hat{U}^n(p_k, q_k) = (p_k, q_k)$$

and M_n is the number of these periodic points for a given n .

Iterating (6.116) n times gives

$$(p_n, q_n) = \begin{cases} (2^n p, \frac{1}{2^n} q) & \text{for } 0 \leq p \leq \frac{1}{2} \\ (2^n p - c_n, \frac{1}{2^n} q + d_n) & \text{for } \frac{1}{2} \leq p \leq 1 \end{cases}$$

where c_n & d_n are some constants independent of p & q .

Therefore,

$$\frac{\partial p_n}{\partial p} = 2^n \quad \frac{\partial q_n}{\partial q} = 2^{-n} \quad \forall q, p$$

(2a) thus becomes,

$$\delta(p - p^n) \delta(q - q^n) = \sum_{k=1}^{M_n} (2^n - 1)^{-1} (1 - 2^{-n}) \delta(q - q_k) \delta(p - p_k)$$

so that (1) gives

$$\text{Tr } \hat{U}^n = M_n (2^n - 1)^{-1} (1 - 2^{-n})^{-1}$$

Now, the points periodic under \hat{U}^n are points with s invariant under n consecutive right bit-shifts. This means s must be of the form

$$s = (\dots r_n r_n r_n \cdot r_n r_n r_n \dots)_2$$

where r_n is a product of n binary digits. Since each digit can have only two possible values, 0 & 1, the possible number of s is $M_n = 2^n$.

Therefore,

$$\begin{aligned} \text{Tr } \hat{U}^n &= 2^n (2^n - 1)^{-1} (1 - 2^{-n})^{-1} \\ &= (1 - 2^{-n})^{-2} \\ &= \sum_{m=0}^{\infty} \frac{(2 + m - 1)!}{(2 - 1)! m!} (2^{-n})^m \\ &= \sum_{m=0}^{\infty} (m + 1) \left(\frac{1}{2^n} \right)^m \end{aligned} \quad (3)$$

By definition, $\text{Tr } \hat{U}^n$ is equal to the sum of the eigenvalues of \hat{U}^n .